

Chapter 6

Programming the Basic Computer – Part 2

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■ 6-5 Program Loops

◆ Program Loops

- A sequence of instructions that are executed many times

◆ Example of program loop

- Sum of 100 integer numbers

» **Fortran**


```
DIMENSION A(100)
INTEGER SUM, A
SUM = 0
DO 3 J = 1, 100
3 SUM = SUM + A(J)
```

Tab. 6-13 Symbolic Program to Add 100 numbers

Line				
1		ORG	100	
2		LDA	ADS	/ A = 150
3		STA	PTR	/ PTR = 150
4		LDA	NBR	/ A = -100
5		STA	CTR	/ CTR = -100
6		CLA		/ A = 0
7	LOP,	ADD	PTR I	/ A + 75
8		ISZ	PTR	/ 150 + 1 = 151
9		ISZ	CTR	/ -100 + 1 = -99
10		BUN	LOP	/ Loop until CTR = 0
11		STA	SUM	/ Store A to SUM
12		HLT		
13	ADS,	HEX	150	
14	PTR,	HEX	0	/ 150
15	NBR,	DEC	-100	
16	CTR,	HEX	0	/ -100
17	SUM,	HEX	0	/ Result of Sum
18		→ ORG	150	
19		DEC	75	
,		,	,	
,		,	,	
118		DEC	23	
119		END		

Data

TABLE 6-13 Symbolic Program to Add 100 Numbers

Line			
1		ORG 100	/Origin of program is HEX 100
2		LDA ADS	/Load first address of operands
3		STA PTR	/Store in pointer
4		LDA NBR	/Load minus 100
5		STA CTR	/Store in counter
6		CLA	/Clear accumulator
7	LOP,	ADD PTR I	/Add an operand to AC
8		ISZ PTR	/Increment pointer
9		ISZ CTR	/Increment counter
10		BUN LOP	/Repeat loop again
11		STA SUM	/Store sum
12		HLT	/Halt
13	ADS,	HEX 150	/First address of operands
14	PTR,	HEX 0	/This location reserved for a pointer
15	NBR,	DEC -100	/Constant to initialize counter
16	CTR,	HEX 0	/This location reserved for a counter
17	SUM,	HEX 0	/Sum is stored here
18		ORG 150	/Origin of operands is HEX 150
19		DEC 75	/First operand
.			
.			
.			
118		DEC 23	/Last operand
119		END	/End of symbolic program

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Program to add two numbers

- Reserve 100 words of memory for **100** operands.
- The numbers are **integers**.
 - If they were of the **float** type,
 - compiler **reserves locations** for floating-point numbers
 - generate **instructions** that perform the subsequent **arithmetic** with **floating-point** data.
- **DIM** and **INTEGER** - **nonexecutable** statements similar assembly pseudoinstructions
- Suppose that the compiler reserves **locations (150)₁₆ to (1B3)₁₆** for the 100 operands.
 - These reserved memory words are listed in **lines 19 to 118**
 - Done by the **ORG pseudoinstruction** in line 18, which specifies the origin of the operands.

Program to add two numbers – cont.

- The **first** and **last** operands are listed with a specific **decimal number**
 - These **values** are **not known** during compilation.
 - Compiler **just reserves** the data space in memory
 - **Values** are **inserted later** when an **input data statement** (not listed in the program)
- Line numbers are for reference only
 - not part of the translated symbolic program.

Program to add two numbers – cont.

- **Line 9:** Only the increment part of ISZ is used
- **AC** is used for **SUM**
 - More efficient than to use a memory location
- **PTR, CTR** are **memory** words
 - When **more registers** are **available (RISC)** an intelligent compiler will **use registers**

■ 6-6 Programming Arithmetic & Logic Operations

◆ Hardware implementation

- Operations are implemented in a computer with one machine instruction
- Ex) **ADD, AND**

◆ Software implementation

- Operations are implemented by a set of instruction(Subroutine)
- Ex) **MUL, DIV**

Hardware - faster and expensive
Software - slower and cheaper

◆ Multiplication Program

- Positive Number Multiplication

» X = multiplicand

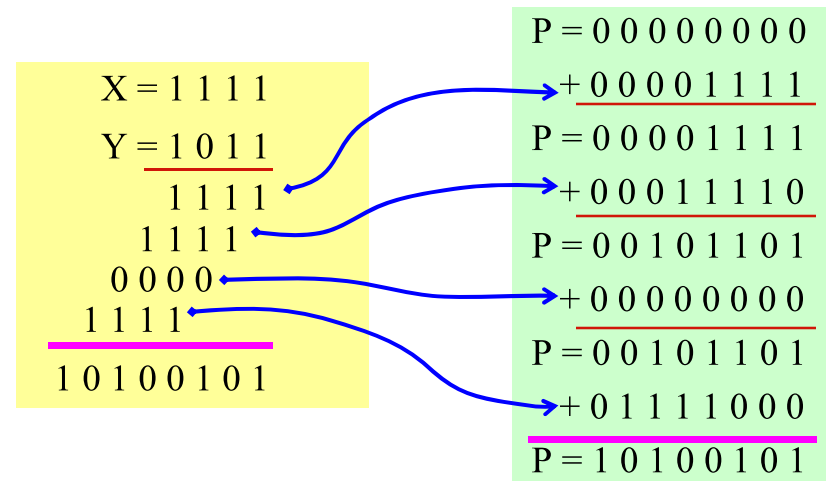
Y = multiplier

P = Partial Product Sum

Algorithm
Fig. 6-3

Circular Right

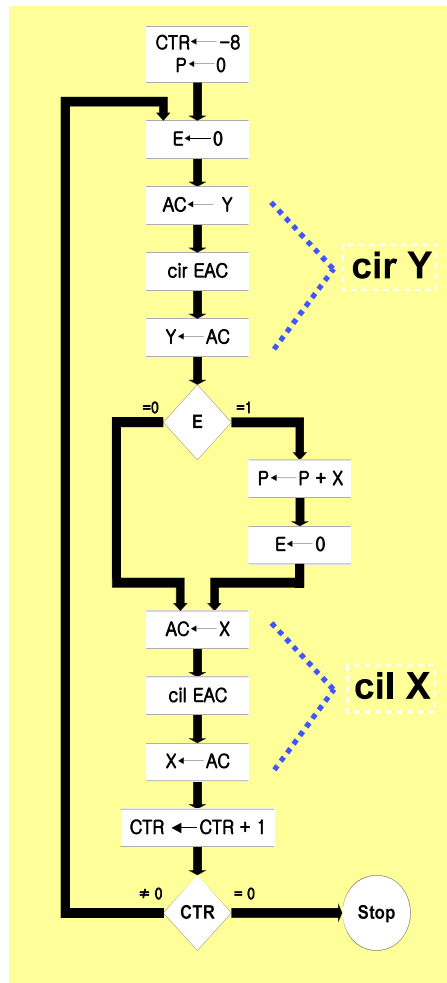
- E = 1
- E = 0



Multiplication program

- **Positive numbers** - disregard sign bit and
- No more than **eight significant** bits
 - their **product** cannot exceed the word capacity of **16 bits**.
 - for **16-bit numbers product** may be up to **31 bits** in length and will occupy two words of memory.
- Solution (like **pen** and **paper**)
 - checking the bits of the **multiplier Y**
 - adding the multiplicand **X** as many times as there are **1's** in **Y**,
 - the value of **X** is shifted left from one line to the next.
- Reserve a memory location, **P**,
 - to store intermediate sums (partial products)
 - since computer can add only two numbers at a time,
 - **P** starts with zero

Fig. 6-3 flowchart for Multiplication Program



Tab. 6-14 Program to Multiply Two Positive numbers

Line				
1		ORG	100	
2	LOP,	CLE		/ A = 0
3		LDA	Y	/ A = Y (000B)
4		CIR		/ Circular Right to E
5		STA	Y	/ Store shifted Y
6		SZE		/ Check if E = 0
7		BUN	ONE	/ E = 1
8		BUN	ZRO	/ E = 0
9	ONE,	LDA	X	A = X (000F)
10		ADD	P	/ X = X + P
11		STA	P	/ St p
12		CLE		/ Clear E
13	ZRO,	LDA	X	/ A = X
14		CIL		/ A = 00011110 (00001111)
15		STA	X	/ St p
16		ISZ	CTR	/ CTR = - 7 = -8 + 1
17		BUN	LOP	/ Repeat until CTR = 0
18		HLT		
19	CTR,	DEC	-8	
20	X,	HEX	000F	
21	Y,	HEX	000B	
22	P,	HEX	0	
23		END		

Alternative?

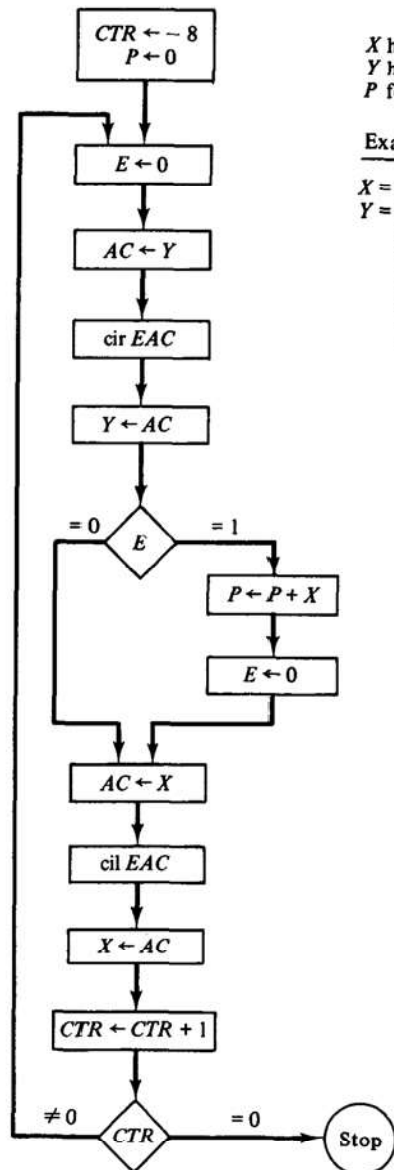


Figure 6-3 Flowchart for multiplication program.

X holds the multiplicand
Y holds the multiplier
P forms the product

Example with four significant digits

<i>X</i> = 0000 1111	<i>P</i>
<i>Y</i> = 0000 1011	0000 0000
0000 1111	0000 1111
0001 1110	0010 1101
0000 0000	0010 1101
0111 1000	1010 0101
1010 0101	

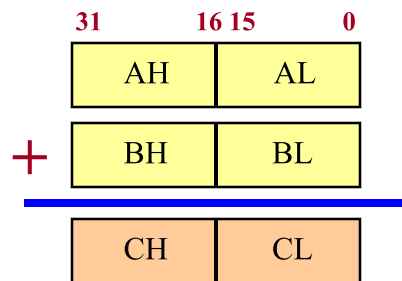
TABLE 6-14 Program to Multiply Two Positive Numbers

	ORG 100	
LOP,	CLE	/Clear <i>E</i>
	LDA Y	/Load multiplier
	CIR	/Transfer multiplier bit to <i>E</i>
	STA Y	/Store shifted multiplier
	SZE	/Check if bit is zero
	BUN ONE	/Bit is one; go to ONE
	BUN ZRO	/Bit is zero; go to ZRO
ONE,	LDA X	/Load multiplicand
	ADD P	/Add to partial product
	STA P	/Store partial product
	CLE	/Clear <i>E</i>
ZRO,	LDA X	/Load multiplicand
	CIL	/Shift left
	STA X	/Store shifted multiplicand
	ISZ CTR	/Increment counter
	BUN LOP	/Counter not zero; repeat loop
	HLT	/Counter is zero; halt
CTR,	DEC -8	/This location serves as a counter
X,	HEX 000F	/Multiplicand stored here
Y,	HEX 000B	/Multiplier stored here
P,	HEX 0	/Product formed here
	END	

◆ Double Precision Addition : 32 bits

• AL + BL

E (AH + BH + E)



Line				
1		LDA	AL	/ A = AL
2		ADD	BL	/ A = AL + BL
3		STA	CL	/ Store A to CL
4		CLA		/ A = 0
5		CIL		/ 0000 0000 0000 000? (?=E)
6		ADD	AH	/ A = 00(E=0) or 01(E=1)
7		ADD	BH	/ A = A + AH + BH
8		STA	CH	/ Store A to CH
9		HLT		
10	AL,	DEC	?	/ Operand
11	AH,	DEC	?	
12	BL,	DEC	?	
13	BH,	DEC	?	
14	CL,	HEX	0	
15	CH,	HEX	0	

◆ Logic Operations

• Logic Operation

OR - How ? DeMorgan's law

$$\gg A + B = \overline{\overline{A} + \overline{B}} = \overline{\overline{A} \cdot \overline{B}}$$



מערכות אוניברסליות - NAND

LDA	A		/ Load A
CMA			/ Complement A
STA	TMP	P	/ St o
LDA	B		/ Load B
CMA			/ Complement
AND	TMP		/ AND
CMA			/ Complement

Program to Add Two Double-Precision Numbers

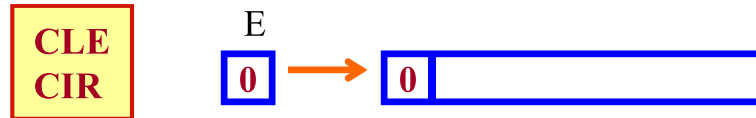
TABLE 6-15 Program to Add Two Double-Precision Numbers

	LDA AL	/Load A low
	ADD BL	/Add B low, carry in E
	STA CL	/Store in C low
	CLA	/Clear AC
	CIL	/Circulate to bring carry into AC(16)
	ADD AH	/Add A high and carry
	ADD BH	/Add B high
	STA CH	/Store in C high
	HLT	
AL,	—	/Location of operands
AH,	—	
BL,	—	
BH,	—	
CL,	—	
CH,	—	

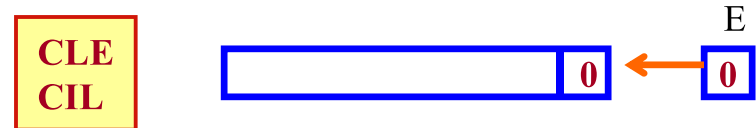
◆ Shift Operations

- Logical Shift : Zero must added to the extreme position

» Shift Right



» Shift Left



- Arithmetic Shift Right

» Positive (+ = 0)



» Negative (- = 1)



CLE	/ E= 0
SPA	/ Skip if A= +, E= 0
CME	/ Toggle E(=1) if A= -
CIR	/ Circulate A with E

■ 6-7 Subroutines

◆ Subroutine

- A set of common instruction that can be used in a program *many times*
- In basic computer, the link between the **main program** and a **subroutine** is the **BSA** instruction(*Branch and Save return Address*)
- Subroutine example : **Tab. 6-16**

Location				
		ORG	100	/ Main Program
100		LDA	X	/ Load X
101		BSA	SH4	/ Call SH4 with X
102		STA	X	/ Store result X
103		LDA	Y	/ Load Y
104		BSA	SH4	/ Call SH4 with Y
105		STA	Y	/ Store result Y
106		HLT		
107	X,	HEX	1234	/ Result = 2340
108	Y,	HEX	4321	/ Result = 3210
				/ Subr
109	SH4,	HEX	0	/ Save Return Address
10A		CIL		
10B		CIL		
10C		CIL		
10D		CIL		
10E		AND	MSK	/ Mask lower 4 bit
10F		BUN	SH4 I	/ Indirect Return to main
110	MSK,	HEX	FFF0	/ Mask pattern
110		END		

Subroutine
CALL hear

X = 102
Y = 105

Tab. 6-16 Program to Demonstrate the use of Subroutines

TABLE 6-16 Program to Demonstrate the Use of Subroutines

Location			
		ORG 100	/Main program
100		LDA X	/Load X
101		BSA SH4	/Branch to subroutine
102		STA X	/Store shifted number
103		LDA Y	/Load Y
104		BSA SH4	/Branch to subroutine again
105		STA Y	/Store shifted number
106		HLT	
107	X,	HEX 1234	
108	Y,	HEX 4321	
			/Subroutine to shift left 4 times
109	SH4,	HEX 0	/Store return address here
10A		CIL	/Circulate left once
10B		CIL	
10C		CIL	
10D		CIL	/Circulate left fourth time
10E		AND MSK	/Set AC(13–16) to zero
10F		BUN SH4 I	/Return to main program
110	MSK,	HEX FFF0	/Mask operand
		END	

◆ Subroutine Parameters & Data Linkage


- Parameter(or Argument) Passing
 - » When a subroutine is called, the main program must transfer the data
- Parameter Passing
 - » 1) Data transfer through the **Accumulator**
 - Used for only single input and single output parameter
 - » 2) Data transfer through the **Memory**
 - Operand are often **placed in memory locations following the CALL**
- 2 Parameter Passing **Tab. 6-17**
 - » **First Operand and Result** : Accumulator
 - » **Second Operand** : Inserted in location following the **BSA**
- BSA 2 Operand : **Tab. 6-18**
 - » BSA 2 Operand
 - » Block Source Destination Address.

Tab. 6-17 Program to Demonstrate Parameter Linkage

Location			
	ORG	200	
200	LDA	X	/ Load first operand X
201	BSA	OR	/ Call OR with X
202	HEX	3AF6	/ Second operand
203	STA	Y	/ Subroutine return here(Y=result)
204	HLT		
205	X, HEX	7B95	/ First operand
206	Y, HEX	0	/ Result store here
207	OR, HEX	0	/ Return address = 202
208	CMA		/ Complement X
209	STA	TMP	/ TMP = X
20A	LDA	OR I	/ A = 3AF6 (202)
20B	CMA		/ Complement Second operand
20C	AND	TMP	/ AND
20D	CMA		/ Complement
20E	ISZ	OR	/ Return Address = 202 + 1 = 203
20F	BUN	OR I	/ Return to main
210	TMP, HEX	0	
	END		

* OR Subroutine
 First Operand : X = 7B95
 Second Operand : BSA = 3AF6

TABLE 6-17 Program to Demonstrate Parameter Linkage

Location			
		ORG 200	
200		LDA X	/Load first operand into AC
201		BSA OR	/Branch to subroutine OR
202		HEX 3AF6	/Second operand stored here
203		STA Y	/Subroutine returns here
204		HLT	
205	X,	HEX 7B95	/First operand stored here
206	Y,	HEX 0	/Result stored here
207	OR,	HEX 0	/Subroutine OR
208		CMA	/Complement first operand
209		STA TMP	/Store in temporary location
20A		LDA OR I	/Load second operand
20B		CMA	/Complement second operand
20C		AND TMP	/AND complemented first operand
20D		CMA	/Complement again to get OR
20E		ISZ OR	/Increment return address
20F		BUN OR I	/Return to main program
210	TMP,	HEX 0	/Temporary storage
		END	

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Tab. 6-18 Subroutine to Move a Block of Data

		ORG	100	
100		BSA	MVE	/ Subroutine Call
101		HEX	200	/ Source Address
102		HEX	300	/ Destin Addressa
103		DEC	-16	/ Number of data to move
104		HLT		
105	MVE,	HEX	0	/ Return address= 101
106		LDA	MVE I	/ A= 200
107		STA	PT1	/ PT1= 200
108		ISZ	MVE	/ Return address= 102
109		LDA	MVE I	/ A= 300
10A		STA	PT2	/ PT2= 300
10B		ISZ	MVE	/ Return address= 103
10C		LDA	MVE I	/ A= -16
10D		STA	CTR	/ CTR= -16
10E		ISZ	MVE	/ Return address= 104
10F	LOP,	LDA	PT1 I	/ A= Address 200
110		STA	PT2 I	/ Address 300
111		ISZ	PT1	/ PT1= 201
112		ISZ	PT2	/ PT2= 301
113		ISZ	CTR	/ CTR= -15 if 0 skip
114		BUN	LOP	/ Loop until CTR= 0
115		BUN	MVE I	/ 104 Retur = HLT
116	PT1,	HEX	?	/ Source
117	PT2,	HEX	?	/ Destination
118	CTR,	DEC	?	/ Counter

2
Operand

Subroutine to Move a **Block** of Data

TABLE 6-18 Subroutine to Move a Block of Data

		/Main program
	BSA MVE	/Branch to subroutine
	HEX 100	/First address of source data
	HEX 200	/First address of destination data
	DEC -16	/Number of items to move
	HLT	
MVE,	HEX 0	/Subroutine MVE
	LDA MVE I	/Bring address of source
	STA PT1	/Store in first pointer
	ISZ MVE	/Increment return address
	LDA MVE I	/Bring address of destination
	STA PT2	/Store in second pointer
	ISZ MVE	/Increment return address
	LDA MVE I	/Bring number of items
	STA CTR	/Store in counter
	ISZ MVE	/Increment return address
LOP,	LDA PT1 I	/Load source item
	STA PT2 I	/Store in destination
	ISZ PT1	/Increment source pointer
	ISZ PT2	/Increment destination pointer
	ISZ CTR	/Increment counter
	BUN LOP	/Repeat 16 times
	BUN MVE I	/Return to main program
PT1,	—	
PT2,	—	
CTR,	—	

3 parameters

■ 6-8 Input-Output Programming

◆ One-character I/O

● Programmed I/O

Tab. 6-19 Program to input and output One character

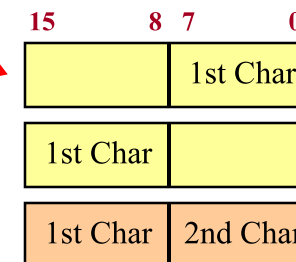
(a) Input a character			
CIF,	SKI		/ Check FGI = 1 ?
	BUN	CIF	/ Go to CIF if FGI= 0
	INP		/ Input character (FGI = 1)
	OUT		/ Echo Back
	STA	CHR	/ Store character
	HLT		
CHR,	?		/ Store character here
(b) Output a character			
	LDA	CHR	/ Load output character
COF,	SKO		/ Check FGO = 1 ?
	BUN	COF	/ Go to COF if FGO= 0
	OUT		/ Output character (FGO = 1)
	HLT		
CHR,	HEX	0057	/ Output character = "W"

◆ Two-character I/O

● Two character Packing

Tab. 6-20 Subroutine to input and pack Two character

IN2,	HEX	?	/ Save return address
FST,	SKI		/ Check if FGI= 1 ?
	BUN	FST	/ Loop (FGI = 0)
	INP		/ Input first character
	OUT		/ Echo back
	BSA	SH4	/ Shift left 4 bit
	BSA	SH4	/ Again(total 8 bit shift)
SCD,	SKI		
	BUN	SCD	
	INP		/ Input second character
	OUT		/ Echo back
	BUN	IN2 I	/ Return



◆ Store Input Character in Buffer

Tab. 6-21 Program to store input character in buffer

	LDA	ADS	/ Load buffer address A= 500
	STA	PTR	/ PTR= 500
LOP,	BSA	IN2	/ Get a character (Tab. 6-20)
	STA	PTR I	/ 500 character
	ISZ	PTR	/ PTR= 501
	BUN	LOP	/ Endless Loop
	HLT		
ADS,	HEX	500	/ Buffer address
PTR,	HEX	0	/ Pointer

◆ Compare Two Word

Tab. 6-22 Program to compare Two word

	LDA	WD1	/ Load first word A= WD1
	CMA		/ Make 2's complement
	INC		
	ADD	WD2	/ WD2 – WD1
	SZA		/ Skip if A=0 (Equal)
	BUN	UEQ	/ Unequal
	BUN	EQL	/ Equal
WD1,	HEX	?	/ first word
WD2,	HEX	?	/ second wor

*Useful for a search
procedure e.g. in look-up
tables*

Remarks

- Can write SH8 instead of double call to SH4
- Program uses a pointer to keep track of current empty location in the buffer.
- No counter is used in the program
- Characters are read
 - as long as they are available or
 - until the buffer reaches location 0 (after location FFFF).
 - In a practical situation - limit the size of the buffer, use a counter

Program to Service an Interrupt

- In former I/O example – **busy waiting**
 - Most running is wasted waiting for external devices to set flags
- Solved by **interrupt facility**
 - notify the computer when a flag is set.
- **Advantage:**
 - **information transfer** only **upon request** from external device.
 - **Meanwhile**, the computer performs other tasks.
- To be **effective**: other program(s) must reside in memory
 - **Multiprogramming environment**

Program to Service an Interrupt – cont.

- **Only one** program can be **executed** at any **given time**
 - However, **two or more** programs **may reside** in memory.
- Program currently being executed - ***running program***.
 - **Other programs** are usually **waiting** for **I/O** data.
- Interrupt facility service procedure
 - **Take care** of the **data transfer** of one (or more) program while another program is currently being executed.
 - The running program must include an **ION** instruction
 - to **turn** the interrupt **on**.
 - **When interrupt** facility is **not used**, program must include an **IOF**

Program to Service an Interrupt – cont.

- Interrupt facility - **allows** the **running program** to **proceed until** the **I/O devices set** their ready flags.
- Whenever a flag is set to 1
 - computer **completes** execution of current **instruction**
 - **Acknowledges** the interrupt.
 - The **return address** is **stored** in location **0**.
 - **Instruction** in **location 1** is **performed** (initiates a **service routine** for the input or output transfer)
- **Service routine** can be **stored anywhere** in memory
 - provided a **branch to the start** of the routine is stored in **location 1**.

Program to Service an Interrupt – cont.

- The service routine must have instructions to perform the following tasks:
 1. **Save contents** of processor registers.
 2. **Check** which **I/O flags** are set.
 3. **Service** the **device** whose flag is set.
 4. **Restore** content of processor **registers**.
 5. **Turn** the **interrupt** facility **on**.
 6. **Return** to the **running program**.
- Also known as a ***Context switching***

Program to Service an Interrupt – cont.

- **Contents** of registers must be **the same**
 - **before** the **interrupt** and **after** the **return** to the running program
 - **otherwise**, the running program may be in **error**
- Service routine may use these **registers**
 - necessary to **save** their **contents** at the beginning of the routine
 - **Restore** them at the **end**.
- Device priority – according to **checking order** of flags
 - higher priority is serviced first, lower served afterwards
- Devices are **serviced one at a time**
 - Although **two or more** flags **may be set** at the **same time**
- During an interrupt other interrupts are ignored
 - Service routine **must turn the interrupt on before returning** to the running program (enable further interrupts)
 - **The interrupt facility should not be turned on until after the return address is inserted into the program counter.**

Program to Service an Interrupt

TABLE 6-23 Program to Service an Interrupt

Location			
0	ZRO,	—	/Return address stored here
1		BUN SRV	/Branch to service routine
100		CLA	/Portion of running program
101		ION	/Turn on interrupt facility
102		LDA X	
103		ADD Y	/Interrupt occurs here
104		STA Z	/Program returns here after interrupt
⋮		⋮	
⋮		⋮	
200	SRV,	STA SAC	/Interrupt service routine
		CIR	/Store content of AC
		STA SE	/Move E into AC(15)
		SKI	/Store content of E
		BUN NXT	/Check input flag
		INP	/Flag is off, check next flag
		OUT	/Flag is on, input character
		STA PT1 I	/Print character
		ISZ PT1	/Store it in input buffer
		NXT,	/Increment input pointer
		SKO	/Check output flag
		BUN EXT	/Flag is off, exit
		LDA PT2 I	/Load character from output buffer
		OUT	/Output character
		ISZ PT2	/Increment output pointer
		LDA SE	/Restore value of AC(15)
		CIL	/Shift it to E
		LDA SAC	/Restore content of AC
		ION	/Turn interrupt on
		BUN ZRO I	/Return to running program
	SAC,	—	/AC is stored here
	SE,	—	/E is stored here
	PT1,	—	/Pointer of input buffer
	PT2,	—	/Pointer of output buffer

Handle
Input

Handle
Output

Store registers

Restore registers

◆ Interrupt Program

● Interrupt Condition

- » Interrupt F/F R = 1
when IEN = 1 and [FGI or FGO = 1]
- » Save return address at 0000
- » Jump to 0001 (Interrupt Start)

● Interrupt Service Routine(ISR)

- » 1) Save Register (AC, E)
- » 2) Check Input or Output Flag
- » 3) Input or Output Service Routine
- » 4) Restore Register (AC, E)
- » 5) Interrupt Enable (ION)
- » 6) Return to the running program

Location				
0	ZR0,	ORG	0	/ Save Interrupt Return Address
1		HEX	?	
		BUN	SRV	/ Jump to ISR
100		ORG	100	/ Main program
101		CLA		
102		ION		/ Turn on Interrupt(IEN= 1)
103		LDA	X	
104		ADD	Y	/ Interrupt occurs here
		STA	Z	/ Return Address(104)
200	SRV,	ORG	200	
201		STA	SAC	/ Save A to SAC
202		CIR		/ Move A into A(15)
203		STA	SE	/ Save
204		SKI		/ Check if FGI= 1?
205		BUN	NXT	/ No, FGI= 0, Check FGO
206		INP		/ Yes, FGI= 1, Character Input
207		OUT		/ Echo back
208		STA	PT1	/ Store in input buffer(PT1)
		ISZ	PT1	/ PT1 + 1
	NXT,	SKO		/ Check if FGO= 1?
		BUN	EXT	/ No, FGO= 0, Exit
		LDA	PT2	/ Yes, FGO= 1, Get output character
		OUT		/ Character output
		ISZ	PT2	/ PT2 + 1
	EXT,	LDA	SE	
		CIL		/ Restore E
		LDA	SAC	/ Restore A
		ION		/ In
		BUN	ZR0	/ Return to running program(104)
	SAC,	HEX	?	
	SE,	HEX	?	
	PT1,	HEX	300	/ Input Buffer Address
	PT2,	HEX	400	/ Output Buffer Address

Interrupt
Here